

Combinatorial Networks  
Week 3, Thursday

Ramsey Number for graphs

- **Definition.** For integer  $k \geq 2$  and integers  $s_1, s_2, \dots, s_k \geq 2$ , the *Ramsey number*  $R_k(s_1, \dots, s_k)$  is the least integer  $n$  such that any  $k$ -edge coloring of  $K_n$  has a clique  $K_{s_i}$  in color  $i$ .
- **Exercise.**  $R_k(s_1, \dots, s_k) < \infty$ .  
(1) (i) If  $k$  is even, then

$$R_k(s_1, \dots, s_k) \leq R_{\frac{k}{2}}(R(s_1, s_2), \dots, R(s_{k-1}, s_k));$$

- (ii) If  $k$  is odd, then

$$R_k(s_1, \dots, s_k) = R_{k+1}(s_1, \dots, s_k, 2).$$

- (2)  $R_k(s_1, \dots, s_k) = \sum_i R_k(s_1, \dots, s_i - 1, \dots, s_k)$ .

- **Exercise.**  $2^k \leq R_k(3, \dots, 3) \leq (k+1)!$ .  
(Remind: For the first inequality, consider bipartite graph).

Ramsey Number for hypergraphs

- **Definition.** Let  $V$  be a finite set,  $2^V = \{A \subset V\}$ ,  $\binom{V}{r} = \{A \subset V : |A| = r\}$ ,  $G = (V, \text{any collection of subsets of } V)$  is called a *hypergraph*, and we call  $H = (V, E)$ , where  $E \subset \binom{V}{r}$  is a  *$r$ -uniform hypergraph*.
- **Note:** 2-uniform hypergraph = graph.
- *complete  $r$ -uniform hypergraph* is  $K_n^{(r)} = (V, \binom{V}{r})$ , where  $|V| = n$ .
- An *independent set* in  $r$ -uniform hypergraph  $H$  is a subset  $S$  of vertices containing NO hyper edge.
- A *clique* in  $H$  is a set of vertices which induced a complete  $r$ -uniform hypergraph.
- **Definition.** The *hypergraph Ramsey number*  $R^{(r)}(s, t)$  is the least  $n$  such that any 2-edge-coloring of  $K_n^{(r)}$  has a blue  $K_s^{(r)}$  or a red  $K_t^{(r)}$ . The Ramsey number  $R^{(r)}(s_1, \dots, s_k)$  is the least  $n$  such that any  $k$ -edge-coloring of  $K_n^{(r)}$  has a monochromatic clique  $K_{s_i}^{(r)}$  (with color  $i$ ).
- **Theorem.** For any  $s, t \geq r$ , the Ramsey number  $R^{(r)}(s, t) < \infty$ . In fact, we prove  $R^{(r)}(s, t) \leq R^{(r-1)}(R^{(r)}(s-1, t), R^{(r)}(s, t-1)) + 1$ .
- **Proof :** By induction on  $r, s, t$ .  
Base case,

$$R^{(2)}(s, t) < \infty, R^{(r)}(s, r) = s < \infty, R^{(r)}(r, t) = t < \infty.$$

Inductive step: Let  $n = R^{(r-1)}(R^{(r)}(s-1, t), R^{(r)}(s, t-1)) + 1$ , consider any 2-edge-coloring of  $K_n^{(r)}$  and an vertex  $v$ ,  $H = (V - \{v\}, \binom{V - \{v\}}{r-1})$  is a complete  $(r-1)$ -uniform hypergraph. Define a 2-edge-coloring on  $H$  by:  $A \in E(H)$  is colored blue if and only if  $A \cup \{v\}$  is blue

in  $r$ -uniform hypergraph  $K_n^{(r)}$ .

Since  $n - 1 = R^{(r-1)}(R^{(r)}(s - 1, t), R^{(r)}(s, t - 1))$ , this 2-edge-coloring on complete  $(r - 1)$ -uniform hypergraph  $H$  has a blue  $K_{R^{(r)}(s-1,t)}^{(r-1)}$  or a red  $K_{R^{(r)}(s,t-1)}^{(r-1)}$ .

Case 1: There is a blue  $K_{R^{(r)}(s-1,t)}^{(r-1)}$  on  $V - \{v\}$ , let  $T = R^{(r)}(s - 1, t)$ . For any  $A \in \binom{T}{r-1}$ ,  $A \cup \{v\}$  is blue in  $K_n^{(r)}$ . Consider complete  $r$ -uniform hypergraph in  $T$  and the restricted 2-edge-coloring of  $K_n^{(r)}$  on  $T$ . Since  $T = R^{(r)}(s - 1, t)$ , this 2-edge-coloring has a blue  $K_{s-1}^{(r)}$  or a red  $K_t^{(r)}$ . In the latter case, we find a red  $K_t^{(r)}$  in  $K_n^{(r)}$ , done; In the former case, we have a blue  $K_{s-1}^{(r)}$ , adding  $v$  to the set, we get a blue  $K_s^{(r)}$ , also done.

Case 2 is similar. ■

- **Theorem.**  $R^{(r)}(s_1, \dots, s_k) < \infty$ .

### An application of 4-uniform hypergraph on geometry

- **Definition.** In Euclidean space, a set  $P \subset R^2$  is *convex*, if for any two vertices  $x, y \in P$ , the line segment connecting  $x$  and  $y$  are also contained in  $P$ .
- **Definition.** The *convex hull* of a set  $P \subset R^2$  is the smallest convex set containing  $P$ .
- **Definition.** A set  $P \subset R^2$  is in a *convex position* if NO points  $x \in P$  is in the convex hull of the other points in  $P$ .
- **Definition.** A set  $P \subset R^2$  is in a *general position* if NO three points are in a line.
- **Fact 1.** Among any 5 points in general position of the plane, there are always 4 points which are in convex position.
- **Proof.** Consider the convex hull of 5 points. If its convex polygon has 4 or 5 points on it, then we are done, so the convex polygon is a  $\triangle abc$ . Consider the other two points  $x, y$  which are inside  $\triangle abc$ . Consider  $xy$ -line, it must intersect two edges of  $\triangle abc$ , say  $ab$  and  $bc$ , then  $x, y, b, c$  are in convex position. ■
- **Fact 2.** For any  $m$  points, if any 4 points of them are in convex position, then these  $m$  points are in convex position.
- **Proof.** Support NOT. There is a point  $a$  is contained in the convex hull of the other  $m - 1$  points. Consider any triangulation of  $(m - 1)$ -polygon, there is a  $\triangle ijk$  contains  $a$ , a contradiction. ■